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G06F 12/08

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G4A AMC

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WO 97/22933 A US 5778432 A US 5524233 A

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(54) Abstract Title

**Invalidating and flushing a predetermined area of cache memory**

(57) A computer system comprises a cache memory with a plurality of cache lines, a storage area to store a data operand, and an execution unit to operate on data elements in the data operand to invalidate a predetermined portion, such as a page in cache memory, of the cache lines in response to receiving a single instruction. The data operand may be a register location (312) containing a portion of a starting address of the cache line in which data is to be invalidated. This portion may include a plurality of most significant bits of the starting address, which is then shifted by a predetermined number of bits by the execution unit to obtain the starting address. The system may set an invalid bit corresponding to the predetermined area of the cache memory. The system may also be used to copy, i.e. flush, a predetermined area of cache memory to a storage area in response to receiving a single instruction. The flushed portion may then be invalidated.

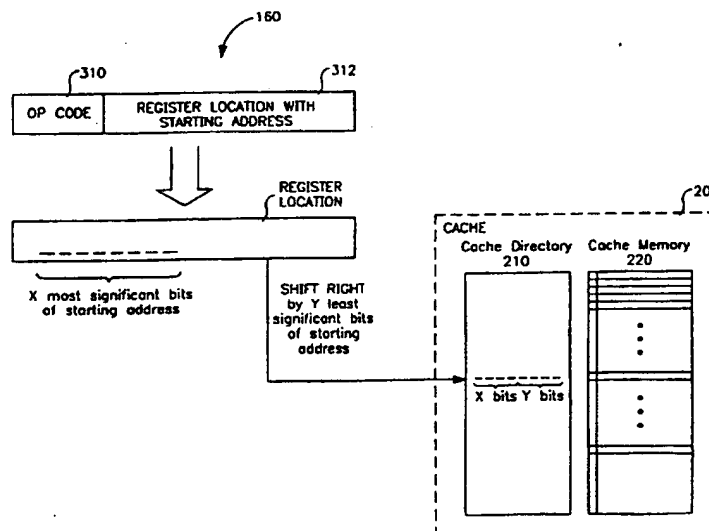


FIG. 3

**GB 2 357 873 A**

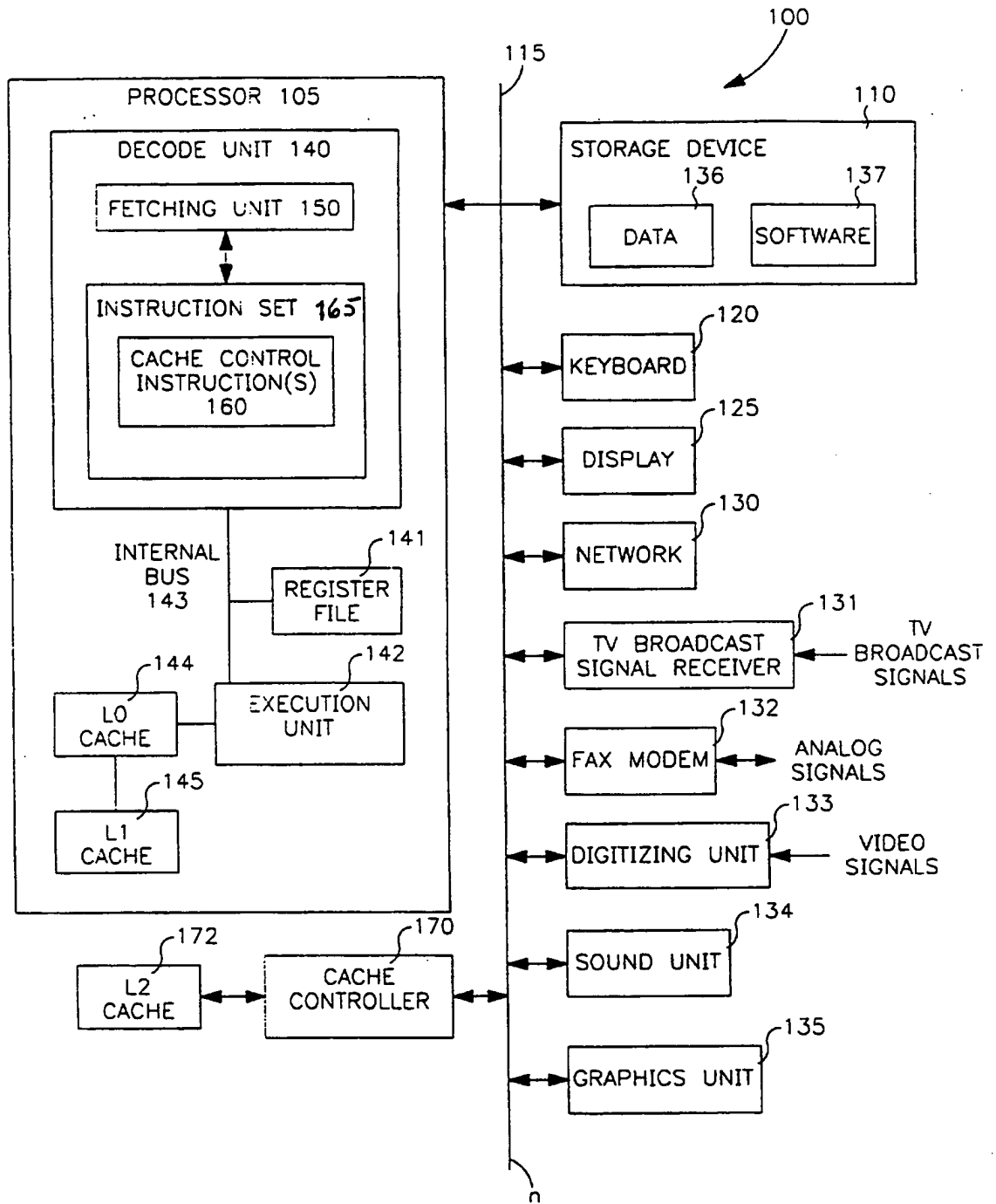


FIG. 1

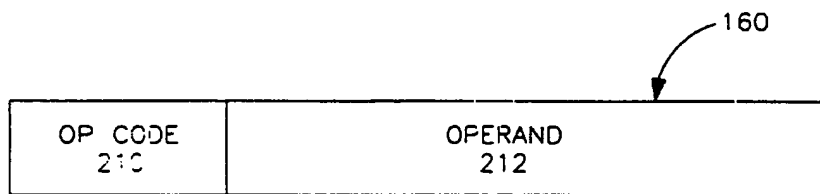


FIG. 2

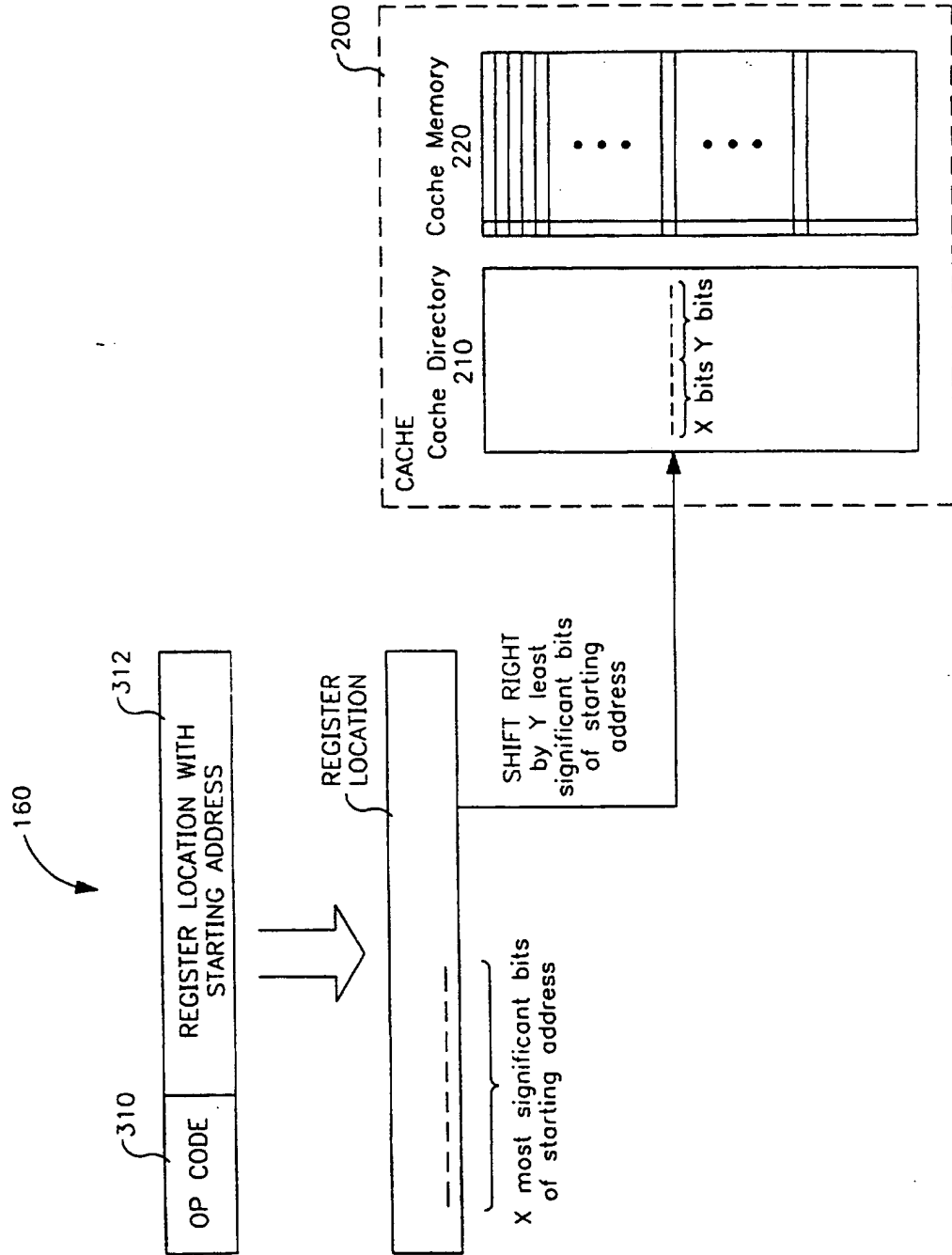


FIG. 3

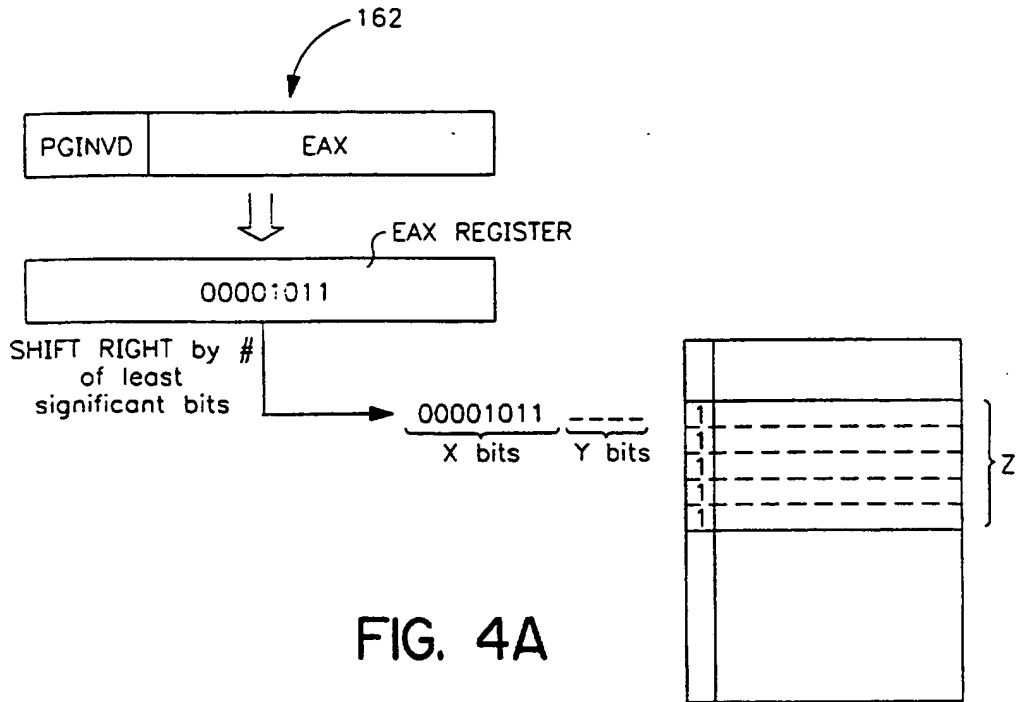


FIG. 4A

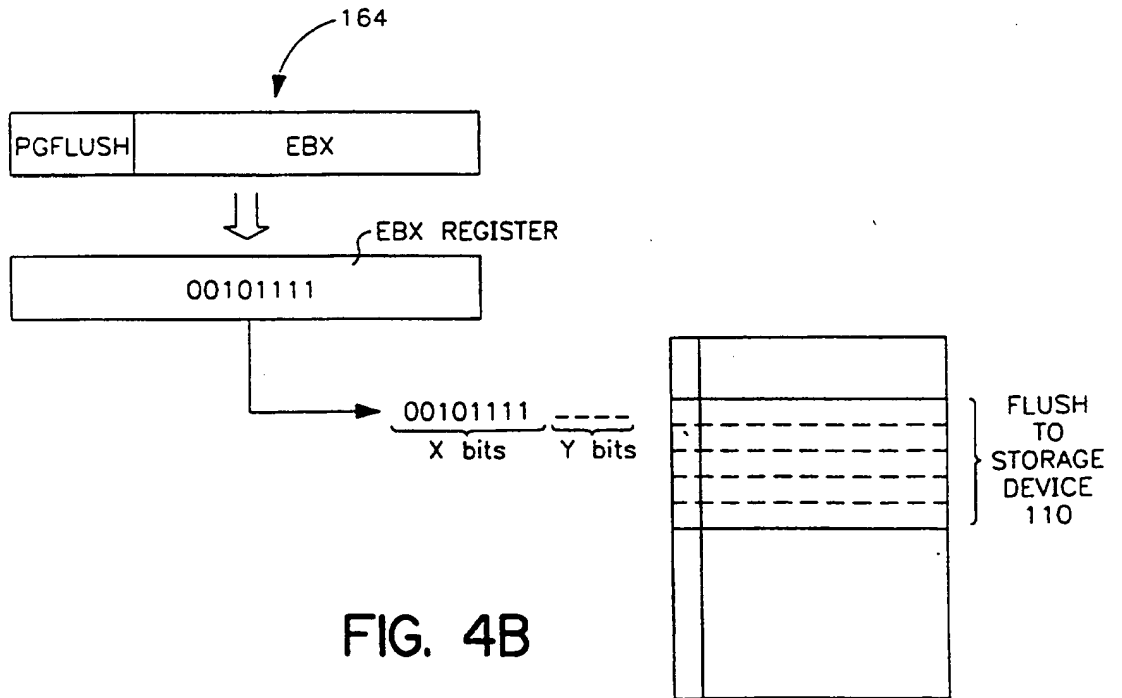


FIG. 4B

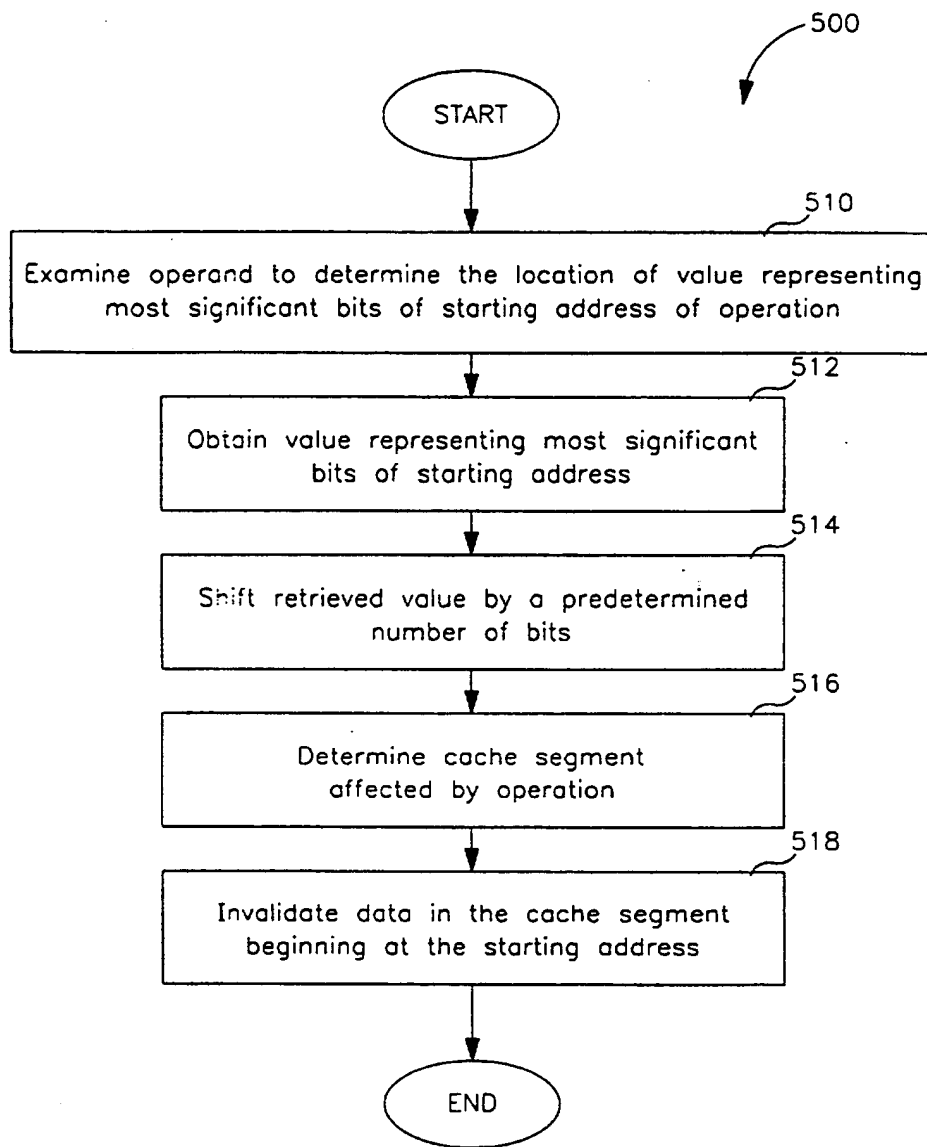


FIG. 5A

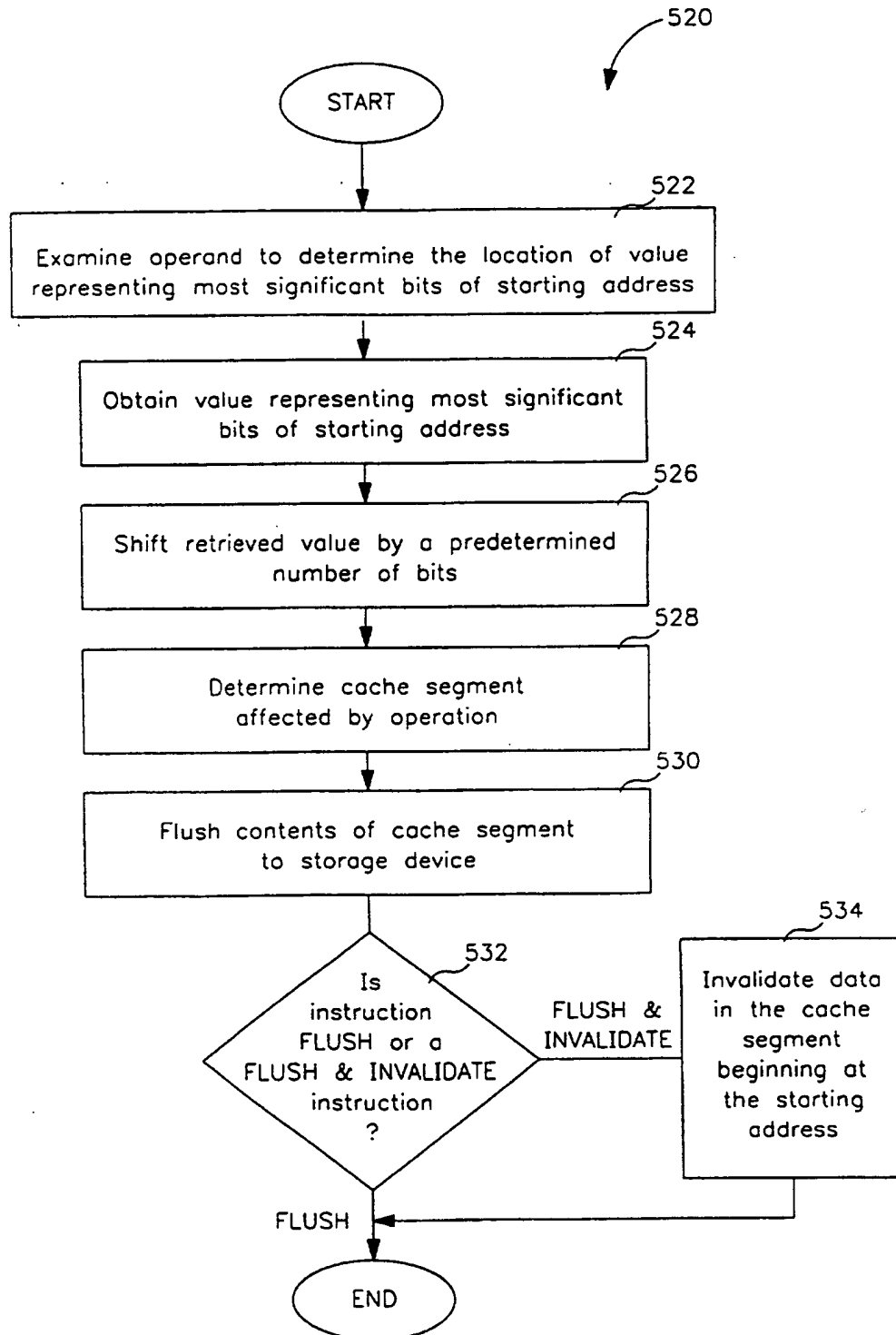


FIG. 5B

A METHOD AND APPARATUS FOR PERFORMING  
CACHE SEGMENT FLUSH OPERATIONS

BACKGROUND OF THE INVENTION

1: Field of the Invention

The present invention relates in general to the field of computer systems, and in particular, to an apparatus and method for providing instructions which facilitate  
5 flushing of a portion of a cache memory within a cache system.

2. Description of the Related Art

The use of a cache memory with a computer system facilitates the reduction of memory access time. The fundamental idea of cache organization is that by  
10 keeping the most frequently accessed instructions and data in the fast cache memory, the average memory access time will approach the access time of the cache. To achieve the optimal tradeoffs between cache size and performance, typical computer systems implement a cache hierarchy, that is, different levels of  
15 cache memory. The different levels of cache correspond to different distances from the computer system core. The closer the cache is to the computer system, the faster the data access. However, the closer the cache is to the computer system, the more costly it is to implement. As a result, the closer the cache level, the faster and smaller the cache.

A cache unit is typically located between the computer system and main  
20 memory; it typically includes a cache controller and a cache memory such as a static random access memory (SRAM). The cache unit can be included on the same chip as the computer system or can exist as a separate component. Alternatively, the cache controller may be included on the computer system chip and the cache memory is formed by external SRAM chips.



The performance of cache memory is frequently measured in terms of its hit ratio. When the computer system refers to memory and finds the data in its cache, it is said to produce a hit. If the data is not found in cache, then it is in main memory and is counted as a miss. If a miss occurs, then an allocation is made at the entry indexed by the address of the access. The access can be for loading data to the computer system or storing data from the computer system to memory. The cached information is retained by the cache memory until it is no longer needed, made invalid or replaced by other data, in which instances the cache entry is de-allocated.

If other computer systems or system components have access to the main memory, as is the case, for example, with a DMA controller, and the main memory can be overwritten, the cache controller must inform the applicable cache that the data stored within the cache is invalid if the data in the main memory changes. Such an operation is known as cache invalidation. If the cache controller implements a write-back strategy and, with a cache hit, only writes data from the computer system to its cache, the cache content must be transferred to the main memory under specific conditions. This applies, for example, when the DMA chip transfers data from the main memory to a peripheral unit, but the current values are only stored in an SRAM cache. This type of operation is known as a cache flush.

Currently, such invalidating and/or flushing operations are performed automatically by hardware, for an associated cache line. In certain situations, software have been developed to invalidate and/or flush the cache memory. Currently, such software techniques involve the use of an instruction which operates on the entire cache memory corresponding to the computer system from which the instruction originated. However, such invalidation and/or flushing operations require a large amount of time to complete, and provides no granularity

or control for the user to invalidate and/or flush specific data or portions of data from the cache, while retaining the other data within the cache memory intact. When a flushing operation operates only on the entire cache memory, it results in inflexibility and impacts system performance. In addition, where a cache  
5 invalidation operation operates only on the entire cache, data corruption may result.

Our co-pending application GB-A-2343029 concerns invalidating and/or flushing a predetermined portion of cache memory and reference is made thereto.

BRIEF SUMMARY OF THE INVENION

According to a first aspect of this invention there is provided a computer system as claimed in claim 1  
5 herein.

According to a second aspect of this invention there is provided a processor as claimed in claim 7 herein.

According to a third aspect of this invention there is provided a computer-implemented method as clamed in  
10 claim 13 herein.

According to a fourth aspect of this invention there is provided a computer-readable apparatus as clamed in  
claim 19 herein.

According to a fifth aspect of this invention there  
15 is provided a computer program according to claim. 21  
herein.

BRIEF DESCRIPTION OF THE DRAWINGS

The invention is illustrated by way of example, and not limitation, in the figures. Like reference indicate similar elements.

Figure 1 illustrates an exemplary computer system in accordance with one  
5 embodiment of the invention.

Figure 2 illustrates one embodiment of the format of a cache control instruction 160 provided according to one embodiment of the invention.

Figure 3 illustrates the general operation of the cache control technique according to one embodiment of the invention.

10 Figure 4A illustrates one embodiment of the operation of the cache segment invalidate instruction 162.

Figure 4B illustrates one embodiment of the operation of the cache segment flush instruction 164.

15 Figure 5A is a flowchart illustrating one embodiment of the cache segment invalidate process of the present invention.

Figure 5B is a flowchart illustrating one embodiment of the cache segment flush process of the present invention.

## DETAILED DESCRIPTION OF THE INVENTION

In the following description, numerous specific details are set forth to provide a thorough understanding of the invention. However, it is understood that the invention may be practiced without these specific details. In other instances, well-known circuits, structures and techniques have not been shown in detail in order not to obscure the invention.

Figure 1 illustrates one embodiment of a computer system 100 which implements the principles of the present invention. Computer system 100 comprises a computer system 105, a storage device 110, and a bus 115. The computer system 105 is coupled to the storage device 110 by the bus 115. The storage device 110 represents one or more mechanisms for storing data. For example, the storage device 110 may include read only memory (ROM), random access memory (RAM), magnetic disk storage mediums, optical storage mediums, flash memory devices and/or other machine readable mediums. In addition, a number of user input/output devices, such as a keyboard 120 and a display 125, are also coupled to the bus 115. The computer system 105 represents a central processing unit of any type of architecture, such as CISC, RISC, VLIW, or hybrid architecture. In addition, the computer system 105 could be implemented on one or more chips. The storage device 110 represents one or more mechanisms for storing data. For example, the storage device 110 may include read only memory (ROM), random access memory (RAM), magnetic disk storage mediums, optical storage mediums, flash memory devices, and/or other machine-readable mediums. The bus 115 represents one or more buses (e.g., AGP, PCI, ISA, X-Bus, VESA, etc.) and bridges (also termed as bus controllers). While this embodiment is described in relation to a single computer

system computer system, the invention could be implemented in a multi-computer system computer system.

In addition to other devices, one or more of a network 130, a TV broadcast signal receiver 131, a fax/modem 132, a digitizing unit 133, a sound unit 134, and a graphics unit 135 may optionally be coupled to bus 115. The network 130 and fax modem 132 represent one or more network connections for transmitting data over a machine readable media (e.g., carrier waves). The digitizing unit 133 represents one or more devices for digitizing images (i.e., a scanner, camera, etc.). The sound unit 134 represents one or more devices for inputting and/or outputting sound (e.g., microphones, speakers, magnetic main memories, etc.). The graphics unit 135 represents one or more devices for generating 3-D images (e.g., graphics card). Figure 1 also illustrates that the storage device 110 has stored therein data 136 and software 137. Data 136 represents data stored in one or more of the formats described herein. Software 137 represents the necessary code for performing any and/or all of the techniques described with reference to Figures 2, and 4-6. Of course, the storage device 110 preferably contains additional software (not shown), which is not necessary to understanding the invention.

Figure 1 additionally illustrates that the computer system 105 includes decode unit 140, a set of registers 141, and execution unit 142, and an internal bus 143 for executing instructions. The computer system 105 further includes two internal cache memories, a level 0 (L0) cache memory which is coupled to the execution unit 142, and a level 1 (L1) cache memory, which is coupled to the L0 cache. An external cache memory, i.e., a level 2 (L2) cache memory 172, is coupled to bus 115 via a cache controller 170. The actual placement of the various cache memories is a design choice or may be dictated by the computer system architecture. Thus, it is

appreciated that the L1 cache could be placed external to the computer system 105. In alternate embodiments, more or less levels of cache (other than L1 and L2) may be implemented. It is appreciated that three levels of cache hierarchy are shown in Figure 1, but there could be more or less cache levels. For example, the present  
5 invention could be practiced where there is only one cache level (L0 only) or where there are only two cache levels (L0 and L1), or where there are four or more cache levels.

Of course, the computer system 105 contains additional circuitry, which is not necessary to understanding the invention. The decode unit 140, registers 141 and  
10 execution unit 142 are coupled together by internal bus 143. The decode unit 140 is used for decoding instructions received by computer system 105 into control signals and/or micro code entry points. In response to these control signals and/or micro code entry points, the execution unit 142 performs the appropriate operations. The decode unit 140 may be implemented using any number of different mechanisms  
15 (e.g., a look-up table, a hardware implementation, a PLA, etc.). While the decoding of the various instructions is represented herein by a series of if/then statements, it is understood that the execution of an instruction does not require a serial processing of these if/then statements. Rather, any mechanism for logically performing this if/then processing is considered to be within the scope of the  
20 implementation of the invention.

The decode unit 140 is shown including a fetching unit 150 which fetches instructions, and an instruction set 165 for performing operations on data. In one embodiment, the instruction set 165 includes a cache control instruction(s) 160 provided  
in accordance with the present invention. In one embodiment, the cache control  
25 instructions include: a cache segment invalidate instruction(s) 162, a cache segment

flush instruction(s) 164 and a cache segment flush and invalidate instruction(s) provided in accordance with the present invention. An example of the cache segment invalidate instruction(s) 162 includes a Page Invalidate (PGINVD) instruction which operates on a user specified linear address and invalidates the 4k  
5 Byte physical page corresponding to the linear address from all levels of the cache hierarchy for all agents in the computer system that are connected to the computer system. An example of the cache segment flush instruction 164 includes a Page Flush (PGFLUSH) instruction 164 that flushes data in the 4 Kbyte physical page corresponding to the linear address on which the operation is performed. An  
10 example of the cache segment flush and invalidate instruction includes a Page Flush/Invalidate (PGFLUSHINV) instruction that first flushes data in the 4 Kbyte physical page corresponding to the linear address on which the operation is performed, and then invalidates the 4 kilobyte physical page corresponding to the linear address. In alternative embodiments, the cache control instruction(s) may  
15 operate on either a user specified linear or physical address and perform the associated invalidate and/or flush operations in accordance with the principles of the invention.

In addition to the cache segment invalidate instruction(s) 162, the cache segment flush instruction(s) 164, and the cache segment flush and invalidate  
20 instruction(s) computer system 105 can include new instructions and/or instructions similar to or the same as those found in existing general purpose computer systems. For example, in one embodiment the computer system 105 supports an instruction set which is compatible with the Intel® Architecture instruction set used by existing computer systems, such as the Pentium®II computer  
25 system. Alternative embodiments of the invention may contain more or less, as well as different instructions and still utilize the teachings of the invention.



The registers 141 represent a storage area on computer system 105 for storing information, such as control/status information, scalar and/or packed integer data, floating point data, etc. It is understood that one aspect of the invention is the described instruction set. According to this aspect of the invention, the storage area  
5 used for storing the data is not critical. The term data processing system is used herein to refer to any machine for processing data, including the computer systems(s) described with reference to Figure 1.

Figure 2 illustrates one embodiment of the format of any one of the cache segment invalidate instruction 162, the cache segment flush instructions 164, and  
10 the cache segment flush and invalidate instruction provided in accordance with the present invention. For discussion purposes, the instructions 162, 164 and will be referred to as the cache control instruction 160. The cache control instruction 160 comprises an operational code (OP CODE) 210 which identifies the operation of the cache control instruction 160 and an operand 212 which specifies the name of a  
15 register or memory location which holds a starting address of the data object that the instruction 160 will be operating on.

Figure 3 illustrates the general operation of the cache control instruction 160 according to one embodiment of the invention. In the practice of the invention, the cache control instruction 160 provides the register (or memory) location which  
20 holds a starting address of the data object that the instruction 160 will be operating on. In one embodiment, the starting address includes X most significant bits, which are stored in the register (or memory) location, and Y least significant bits. The cache control process associated with the cache control instruction 160 then shifts the X bits to the right by Y bit positions to obtain the complete starting address. The  
25 cache control instruction 160 then operates on the data corresponding to the starting

address, and data corresponding to the Z subsequent addresses, in cache memory. In one embodiment, the cache control instruction 160 operates on one page of data stored in cache, of which the beginning address is stored in a register (or memory) location specified in the operand 212 of the cache control instruction. In alternate  
5       embodiments, the cache control instruction 160 may operate on any predetermined amount of data stored in cache, of which the beginning address is stored in a register (or memory) location specified in the operand 212 of the cache control instruction.

In Figure 1, only L0, L1 and L2 levels are shown, but it is appreciated that more or less levels can be readily implemented. The embodiment shown in Figures  
10       4-6 describes the use of the invention with respect to one cache level.

Details of various embodiments of the cache control instruction 160 will now be described. The cache segment invalidate instruction 162 will first be described. Figure 4A illustrates one embodiment of the cache segment invalidate instruction 162. Upon receiving the cache segment invalidate instruction 162, the computer  
15       system 105 determines, from the operand 312 of the instruction 162, the register location in which the most significant bits of the starting address of the data object is stored. The computer system 105 then shifts the value in the operand 312, by the number of least significant bits of the starting address. Once the complete starting address is obtained, the computer system 105 sets the invalidate bit of the cache  
20       memory 200 corresponding to the affected locations of the cache memory. In one embodiment, one page of the cache memory 220 having a starting address corresponding to that stored in the operand 312 will be invalidated. In alternate embodiments, data in any predetermined portions of the cache memory 220 having a starting address corresponding to that stored in the operand 312 will be invalidated  
25       using the present technique.

The cache segment flush instruction 164 will next be described. Figure 4B illustrates one embodiment of the cache segment flush instruction 164. Upon receiving the cache segment flush instruction 164, the computer system 105 determines, from the operand 312 of the instruction 164, the register location in which the most significant bits of the starting address of the data object is stored. The computer system 105 then shifts the value in the operand 312, by the number of least significant bits of the starting address. Once the complete starting address is obtained, the computer system flushes the locations of cache memory 220 affected by execution of the instruction 164. In one embodiment, one page of the cache memory 220 having a starting address corresponding to that stored in the operand 312 will be flushed. In alternate embodiments, data in any predetermined portions of the cache memory 220 having a starting address corresponding to that stored in the operand 312 will be flushed.

Upon receiving the cache segment flush instruction, the computer system 105 determines, from the operand 312 of the instruction 164, the register location in which the most significant bits of the starting address of the data object is stored. The computer system 105 then shifts the value in the operand 312, by the number of least significant bits of the starting address. Once the complete starting address is obtained, the computer system flushes the locations of cache memory 220 affected by execution of the instruction 164. In one embodiment, one page of the cache memory 220 having a starting address corresponding to that stored in the operand 312 will be flushed. In alternate embodiments, any predetermined portions of the cache memory 220 having a starting address corresponding to that stored in the operand 312 will be flushed. Next, the computer system 105

invalidates the affected areas of the cache memory 220 that have been flushed. In one embodiment, this is performed by setting the invalidate bit of each affected cache line.

Figure 5A is a flowchart illustrating one embodiment of the cache segment  
5 invalidate process of the present invention. Beginning from a start state, the process  
500 proceeds to process block 510, where it examines the operand 312 of the  
instruction 62 received by the computer system 105 to determine the storage location  
of the value representing the most significant bits of the starting address of the  
corresponding operation. The process 500 then proceeds to process block 512, where  
10 it retrieves the value representing the most significant bits of the starting address  
from the storage location specified. The process 500 then advances to process block  
514, where it shifts the retrieved value by a predetermined number of bits. In one  
embodiment, the predetermined number represents the number of least significant  
bits in the starting address. Next, the process 500 determines the cache segment  
15 affected by the operation or the instruction 162, as shown in process block 516. In  
one embodiment, the cache segment is a page. In one embodiment, a page contains  
4k Bytes. In alternate embodiments, the cache segment may be any predetermined  
portion of the cache memory. The process 500 then proceeds to process block 516,  
where it invalidates the data in the corresponding cache segment beginning at the  
20 starting address specified. In one embodiment, this is performed by setting the  
invalid bit corresponding to each cache line in the cache segment. The process 500  
then terminates.

Figure 5B is a flowchart illustrating one embodiment of the cache segment  
flush process of the present invention. Beginning from a start state, the process 520  
25 proceeds to process block 522, where it examines the operand 312 of the instruction

64 or 66 received by the computer system 105 to determine the storage location of the value representing the most significant bits of the starting address of the corresponding operation. The process 520 then proceeds to process block 524, where it retrieves the value representing the most significant bits of the starting address  
5 from the storage location specified. The process 520 then advances to process block 526, where it shifts the retrieved value by a predetermined number of bits. In one embodiment, the predetermined number represents the number of least significant bits in the starting address. Next, the process 520 determines the cache segment affected by the operation or the instruction 64 or 66, as shown in process block 528.  
10 In one embodiment, the cache segment is a page. In alternate embodiments the cache segment may be any predetermined portion of the cache. The process 520 then proceeds to process block 530, where it flushes the contents of the cache segment to the storage device specified. The process 520 then proceeds to decision block 530, where it queries if the instruction received corresponding to the operation is a  
15 FLUSH or a FLUSH and INVALIDATE instruction. If the instruction is a FLUSH, the process 520 terminates. If the instruction is a FLUSH and INVALIDATE instruction, the process 520 proceeds to process block 534, where it invalidates the data in the corresponding cache segment beginning at the starting address specified. In one embodiment, this is performed by setting the invalid bit corresponding to  
20 each cache line in the cache segment. The process 520 then terminates.

The use of the present invention thus enhances system performance by providing an invalidate instruction and/or a flush instruction for invalidating and/or flushing data in any predetermined portion of the cache memory. For cases where consistency between the cache and main memory are maintained by software,  
25 system performance is enhanced, since flushing only the affected portions of cache is more efficient and flexible than flushing the entire cache. In addition, system

performance is enhanced by having a flushing and/or invalidate operation that has a granularity that is larger than a cache line size, since the user can flush and/or invalidate a memory region using a single instruction instead of having to alter the code, as the computer system changes the size of a cache line.

5           While a preferred embodiment has been described, it is to be understood that the invention is not limited to such use. In addition, while the invention has been described in terms of several embodiments, those skilled in the art will recognize that the invention is not limited to the embodiments described. The method and apparatus of the invention can be practiced with modification and alteration within  
10 the scope of the appended claims.       The description is thus to be regarded as illustrative instead of limiting on the invention.

CLAIMS:

1. A computer system comprising:
  - a first storage area to store data;
  - 5 a cache memory having a plurality of cache lines each of which stores data;
  - a second storage area to store a data operand; and
  - an execution unit coupled to said first storage area, said second storage area, and said cache memory, said
  - 10 execution unit to operate on data elements in said data operand to copy data from a predetermined portion of the plurality of cache lines in the cache memory to the first storage area, in response to receiving a single instruction.
2. The computer system of claim 1, wherein the data operand
- 15 is a register location.
3. The computer system of claim 2, wherein the register location contains a plurality of most significant bits of a starting address of the cache line in which data is to be copied.
- 20 4. The computer system of claim 3, wherein execution unit shifts the data elements by a predetermined number of bit positions to obtain the starting address of the cache line in which data is to be copied.
5. The computer system of claim 1, wherein the predetermined
- 25 portion of the plurality of cache lines is a page in the cache memory.
6. The computer system of claim 1, wherein the execution unit further invalidates data in the predetermined portion of the plurality of cache lines in response to receiving the
- 30 single instruction, upon copying the data to the first storage area.

7. A processor comprising:  
a decoder configured to decode instructions, and  
5 a circuit coupled to said decoder, said circuit in response to  
a single decoded instruction being configured to:  
obtain a starting address of a predetermined area of a  
cache memory on which the instruction will be performed;  
copy data in the predetermined area of cache memory;  
10 store the copied data in a storage area separate from the  
cache memory.
8. The processor of claim 7, wherein a portion of the  
starting address is located in a register specified in the  
decoded instruction.
- 15 9. The processor of claim 7, wherein the portion of the  
starting address includes a plurality of most significant bits  
of the starting address.
10. The processor of claim 9, wherein the circuit shifts the  
data elements by a predetermined number of bit positions to  
20 obtain the starting address of the cache line in which data is  
to be copied.
11. The processor of claim 9, wherein the predetermined  
portion of the plurality of cache lines is a page in the cache  
memory.
- 25 12. The processor of claim 9, wherein said circuit further  
invalidates the data in the predetermined portion of the  
plurality of cache lines in response to receiving the single  
instruction, upon copying the data to the storage area.
13. A computer-implemented method, comprising:  
30 a) decoding a single instruction;  
b) in response to said step of decoding the single



instruction, obtaining a starting address of a predetermined area of a cache memory on which the single instruction will be performed; and

5 c) completing execution of said single instruction by copying data in a predetermined area of cache memory and storing the copied data in a storage area separate from the cache memory.

10 14. The method of claim 13, wherein c) comprises setting an invalid bit corresponding to the predetermined area of cache memory.

15 15. The method of claim 13, wherein b) comprises:

b.1) obtaining a portion of the starting address from a storage location specified in the decoded instruction;

b.2) shifting the portion of the starting address by a predetermined number of bit positions to obtain the starting address of the cache line in which data is to be invalidated.

16 16. The method of claim 15, wherein in b.1) the portion of the starting address contains a plurality of most significant bits of the starting address, and wherein in b.2) the predetermined number of bit positions represent the number of least significant bits of the starting address.

17 17. The method of claim 13, wherein the predetermined portion of the plurality of cache lines is a page in the cache memory.

18 18. The method of claim 13, further comprising:

d) invalidating the data in the predetermined portion of the plurality of cache lines in response to receiving the single instruction, upon copying the data to the storage area.

19 19. A computer-readable apparatus comprising:

a computer-readable medium that stores an instruction

which when executed by a processor causes said processor to:  
obtain a starting address of a predetermined area of a  
5 cache memory on which the instruction will be performed;  
copy data from the predetermined area of cache memory;  
and

store the copied data in a storage area separate from the  
cache memory.

10 20. The apparatus of claim 19, wherein the instruction  
further causes the processor to:

invalidate the data in the predetermined portion of the  
plurality of cache lines in response to receiving the  
instruction, upon copying the data to the storage area.

15 21. A computer program comprising computer program code means  
adapted to perform all the steps of claim 13 when that program  
is run on a computer.



INVESTOR IN PEOPLE

Application No: GB 0105382.6  
Claims searched: 1-21

20

Examiner: Ben Micklewright  
Date of search: 25 April 2001

### Patents Act 1977 Search Report under Section 17

#### Databases searched:

UK Patent Office collections, including GB, EP, WO & US patent specifications, in:
UK Cl (Ed.S): G4A (AMC)
Int Cl (Ed.7): G06F (12/08)
Other: Online: WPI, EPODOC, PAJ, INSPEC, COMPUTER

#### Documents considered to be relevant:

Category	Identity of document and relevant passage	Relevant to claims
X	GB 2210480 A (SUN) See whole document, e.g. pages 1,2 and pages 10-17	1-21
X	EP 0817081 A2 (SUN) See e.g. column 6 lines 3-16 and column 11 lines 40-48	1-21
X	EP 0210384 A1 (HEWLETT-PACKARD) See e.g. pages 5-9	1-21
X	WO 97/22933 A1 (AMD) See e.g. pages 3 and 7	1-21
X	US 5778432 (RUBIN) See e.g. column 2 lines 26-37	1-21
X	US 5524233 (MILBURN) See e.g. column 10 lines 10-15	1-21

X Document indicating lack of novelty or inventive step	A Document indicating technological background and/or state of the art.
Y Document indicating lack of inventive step if combined with one or more other documents of same category.	P Document published on or after the declared priority date but before the filing date of this invention.
& Member of the same patent family	E Patent document published on or after, but with priority date earlier than, the filing date of this application.